

# Function Maximization with Dynamic Quantum Search

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## The problem

$$\max_{x \in \{0/1\}^n} f(x)$$



$n = 4$  Knapsack example:

- Item 1 (7 kg, \$40)
- Item 2 (4 kg, \$100)
- Item 3 (2 kg, \$50)
- Item 4 (3 kg, \$30)

# Knapsack 0/1 example

Table: Candidate solutions and values for Knapsack problem data.

Candidate	Fitness / 10	Weight	Validity
0 0 0 0	0	0	valid
0 0 0 1	3	3	valid
0 0 1 0	5	2	valid
0 0 1 1	8	5	valid
0 1 0 0	10	4	valid
0 1 0 1	13	7	valid
0 1 1 0	15	6	valid
<b>0 1 1 1</b>	<b>18</b>	<b>9</b>	<b>valid</b>
1 0 0 0	4	7	valid
1 0 0 1	7	10	valid
1 0 1 0	9	9	valid
1 0 1 1	12	12	invalid
1 1 0 0	14	11	invalid
1 1 0 1	17	14	invalid
1 1 1 0	19	13	invalid
1 1 1 1	22	16	invalid

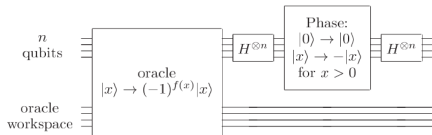
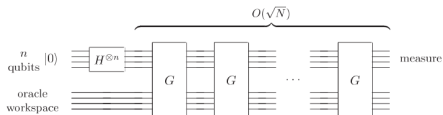
## Interesting quantum properties

- Superposition :  $\sum_{x=0}^{2^n-1} a_x|x\rangle$
- Parallel application :  $\sum_{x=0}^{2^n-1} a_x|x\rangle|0\rangle \rightarrow \sum_{x=0}^{2^n-1} a_x|x\rangle|f(x)\rangle$
- Quadratic speedup with Quantum Search (Grover's algorithm - 1996)
- Adaptation for finding the maximum (Ahuja, Kapoor - 1999)

# Quantum Search

Quantum search in literature with a list  $P$  of  $N = 2^n$  elements:

- Start with :  $\frac{1}{\sqrt{N}} \sum_{x=0}^{N-1} |x\rangle$
- Specify an oracle :  $f(x) = 1$  if  $x$  is the sought element, 0 otherwise
- Apply Grover iterations  $O(\sqrt{N})$  times
- Measure with a high probability to get the sought  $x$ .



## Quantum Search of the Maximum

- For finding the maximum, same process but input an index  $k$  for the maximum :

$$\frac{1}{\sqrt{N}} \sum_{x=0}^{N-1} |x\rangle|k\rangle$$

- Specify an oracle :  $f(x) = 1$  if  $P[x] > P[k]$ , 0 otherwise
- Apply iterations and measure a new  $k$
- Complexity  $\approx 13.6\sqrt{N}$

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## Initial requirements

### Specify registers :

- $q$ : an  $n$ -qubit register for storing candidate solutions.  $n = 4$  in this example.
- $f$ : a  $p$ -qubit register for storing the value of the fitness calculation in two's complement representation, where  $p$  is sufficient to store the sum of all possible values.  $p = 6$  in this example.
- $w$ : a register for storing the total mass of a candidate backpack.  $\dim(w) = 5$  using unsigned integer representation in this example.
- $g$ : a register for storing intermediate computational states and numerical constants.  $\dim(g) = 6$  for this example.
- $v$ : a 1-qubit register to mark a candidate fitness as valid or invalid.
- $r$ : a 1-qubit register used by the oracle.

### Arithmetic quantum circuits for the oracle :

- Let  $a, b$  be integers in  $m$ -bitstring format  $\{a_i, b_i\}_{i=0}^{m-1}$ ,  $a = \sum_{i=0}^{m-1} a_i 2^i$
- Quantum ripple adder :  $|a, b\rangle \rightarrow |a, a + b\rangle$
- Controlled version :  $|c, a, b\rangle \rightarrow |c = 0, a, b\rangle, |c = 1, a, a + b\rangle$
- Comparator :  $|a, b, 0\rangle \rightarrow |a, b, 1\rangle$  if  $a < b$

## Building the oracle : Program

- 1  $\frac{1}{\sqrt{16}} \sum_{i=0}^{15} |i\rangle_q |0\rangle_f |0\rangle_w |0\rangle_g |0\rangle_v |-\rangle_r$
- 2 Control-add  $w_1 : (\sum_{i,i_0=0} |i\rangle_q |0\rangle_w + \sum_{i,i_0=1} |i\rangle |0 + w_1\rangle_w) |w_1\rangle_g$
- 3 Control-add  $w_2, w_3, w_4 : \sum_{i=0}^{15} |i\rangle_q |0\rangle_f |m(i)\rangle_w$
- 4 Control-add  $f_1, f_2, f_3, f_4 : \sum_{i=0}^{15} |i\rangle_q |f(i)\rangle_f |m(i)\rangle_w$
- 5 Define validity with comparator :

$$\sum_{i,m(i)\leq 10} |i\rangle_q |m(i)\rangle_w |10\rangle_g |0\rangle_v + \sum_{i,m(i)>10} |i\rangle_q |m(i)\rangle_w |10\rangle_g |1\rangle_v$$

- 6 Inverse fitness for invalid candidates (by complementing and adding 1):

$$\sum_{i,m(i)\leq 10} |i\rangle_q |f(i)\rangle_f |m(i)\rangle_w |0\rangle_g |0\rangle_v + \sum_{i,m(i)>10} |-f(i)\rangle_f |i\rangle_q |m(i)\rangle_w |0\rangle_g |1\rangle_v$$

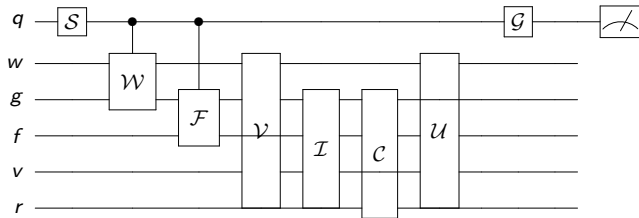
- 7 Mark better candidates using a comparator :

$$\sum_{i,f(i)\leq k} |i\rangle_q |f(i)\rangle_f |k\rangle_g |-\rangle_r + \sum_{i,f(i)>k} -|i\rangle_q |f(i)\rangle_f |k\rangle_g |-\rangle_r$$

- 8 Uncompute :

$$\sum_{i,f(i)\leq k} |i\rangle_q |0\rangle_f |0\rangle_w |0\rangle_g |0\rangle_v |-\rangle_r + \sum_{i,f(i)>k} -|i\rangle_q |0\rangle_f |0\rangle_w |0\rangle_g |0\rangle_v |-\rangle_r$$

## Circuit and Strategy



### Quantum search for the case $M$ is unknown (Boyer et al., 1999)

- 1 Initialize  $m = 1$  and set a constant  $\lambda = 6/5$ .
- 2 Choose a random integer  $j$  that is less than  $m$ .
- 3 Apply  $j$  Grover routines over the superposition of candidates
- 4 Measure to get an outcome.
- 5 Update current max if better.
- 6 Else set  $m$  to  $\min(\lambda m, \sqrt{N})$  and repeat from step 2.

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## Simulation and Discussion



- Could get solution in 2-3 Grover iterations
- Need uncomputation and reversibility
- Alternative to or combine with standard classical techniques

## Publication

*Moussa C., Calandra H., Humble T.S. (2019) Function Maximization with Dynamic Quantum Search. In: Feld S., Linnhoff-Popien C. (eds) Quantum Technology and Optimization Problems. QTOP 2019. Lecture Notes in Computer Science, vol 11413. Springer, Cham*

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<https://arxiv.org/abs/1902.00445>

*Thank you!*